Introduction to Android

CS 436 Software Development on Mobile



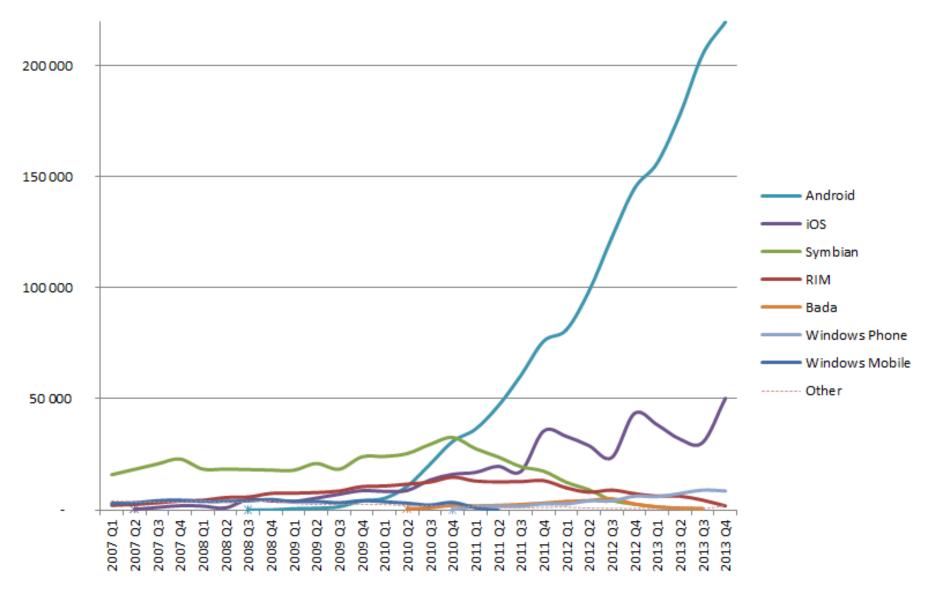


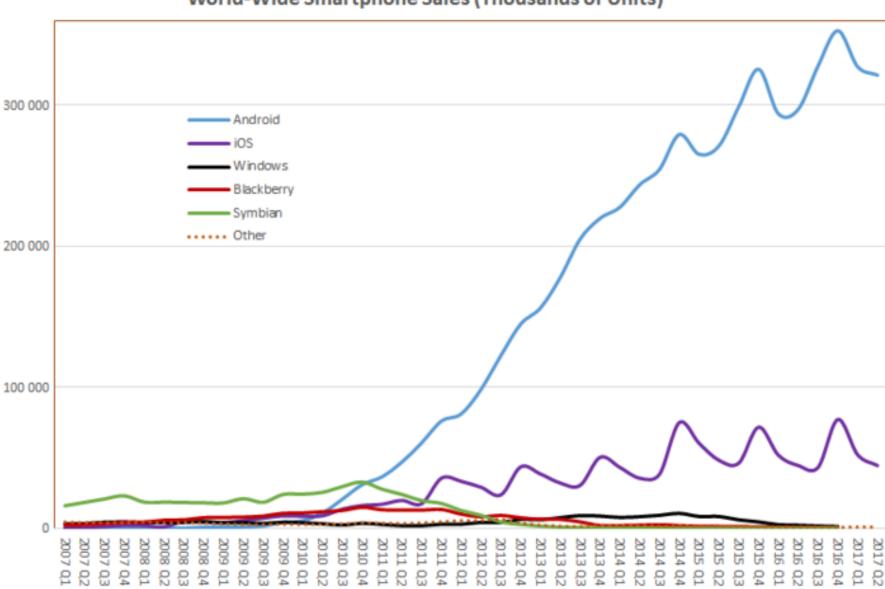
By Dr. Paween Khoenkaw





World-Wide Smartphone Sales (Thousands of Units)





World-Wide Smartphone Sales (Thousands of Units)

What is Android ?

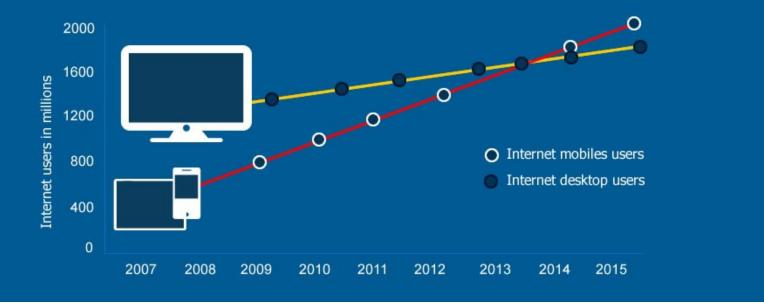
- Complete embedded operating system
- Cutting-edge mobile user experience
- Software stack for building application
- Open platform

Why Android was created ?

- Full phone software stack including application
- Designed as a platform for software development
- Open
- Free
- Community support
- Java Phone

Internet usage - Mobiles VS. Ordinateurs

The projection of global internet users conducted by Morgan Stanley Research: Mobiles VS. Computers from 2007 to 2015.



Android inc.

-Android, Inc. was founded in October 2003 by Andy Rubin.

-The early intentions were to develop an operating system for digital cameras.

-The company diverted its efforts toward producing a smartphone operating system.

-Operating system that would rival Symbian and Microsoft Windows Mobile



Andrew E. "Andy" Rubin

Handset Manufacturers





<u>Mobile Operators</u>









NUANCE



pv

Semiconductor



<u>Commercialization</u>





WIND RIVER

- Software platform from Google and the Open Handset Alliance
- August 2005, Google acquired Android, Inc.
- November 2007, Open Handset Alliance formed to develop open standards for mobile devices
- October 2008, Android available as open source
- December 2008, 14 new members joined Android project



Eric Schmidt, Andy Rubin, and Hugo Barra

Android Versions

- Android Open Source Project (AOSP)
 - Open Source



Firefox OS



Kindle OS

- Google Apps API
 - Closed Source
 - Google Play Service



LINEAGE



2008 Aug 18	ton 2 to 100 to	2009 Feb 9	2009 Apr 30
Android 0.9 320x480 HVGA	Android 1.0 Apple pie API 1.0 320x480 HVGA 320x480 HVGA	Android 1.1 Banana bread API 2.0 320x480 HVGA	Android 1.5 Cupcake API 3.0 Bluetooth A2DP, AVRCP support Soft-keyboard with text- prediction Record/watch videos 320x480 HVGA

2009 Sep 15	2009 Oct 26	2010 May 20
Android 1.6	Android 2.0	Android 2.2
<u>Donut</u>	Éclair	Froyo
<u>API 4.0</u>	API 5.0	API 8.0
Gesture framework	Digital zoom	Flash 10.1
Turn-by-turn navigation	Live Wallpapers	JIT implementation
800×480 WVGA	Updated UI	USB Tethering
	800×480 WVGA	Applications installation to the expandable memory
		Upload file support in the
		browser
		Animated GIFs
		800×480 WVGA



2010 Dec 6



2011 Feb 22

Android 2.3 Gingerbread API 9.0

Near Field Communication support Native VoIP/SIP support Video call support 1366×768 WXGA Android 3.0 Honeycomb API 11.0 Multi core support Better tablet support Updated 3D UI "Private browsing" Open Accessory API USB host API Mice, joysticks, gamepads... support



2011 Oct 19



Android 4.0 Ice Cream Sandwich

<u>API 14.0</u>

Facial recognition (Face Unlock)

UI use Hardware acceleration Web browser, allows up to 16 tabs

Updated launcher

Android Beam exchange data through NFC

Resizable widgets

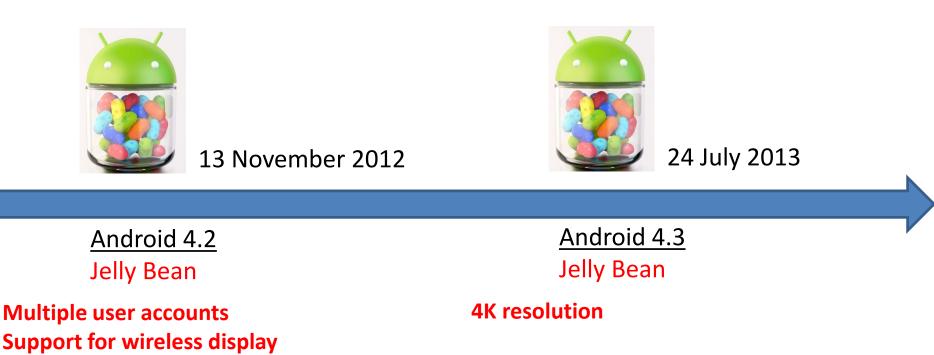
Video stabilization

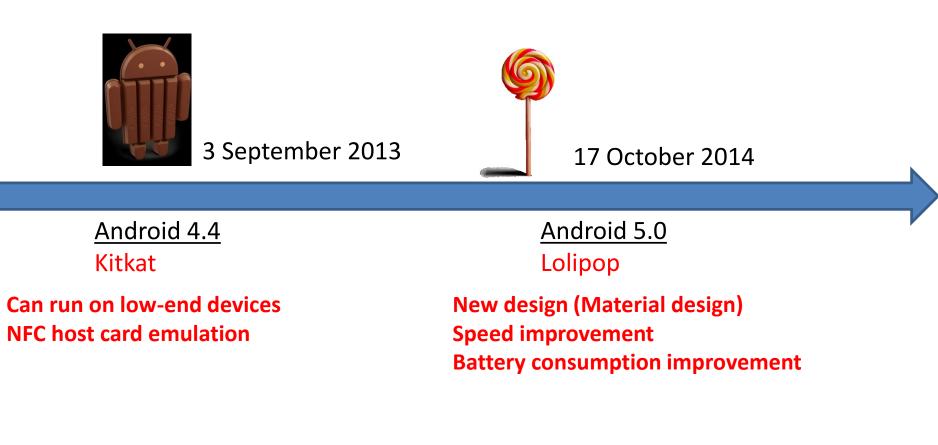
GoogleNow

Android 4.1 Jelly Bean

Triple buffering in the graphics pipeline

Extends vsync timing across all drawing and animation CPU input boost Bi-Directional Text and Other Language Support Android Beam Google Cloud Messaging for Android App Encryption Smart App Updates









22 August 2016

<u>Android 6</u> Marshmallow

5 October 2015

USB Type-C support Fingerprint Authentication support Better battery life with "deep sleep" Permissions dashboard Android Pay MIDI support <u>Android 7</u> Nougat

Unicode 9.0 emoji Better multitasking Multi-window mode (PIP, Freeform window)

Seamless system updates (with dual system partition) Better performance and code size thanks to new JIT Compiler



21 August 2017

December 5, 201

<u>Android 8.0</u> Oreo

Multi-display support 2 times faster boot time Downloadable fonts Integrated printing support Android 8.1

Neuron Network API Shared Memory Android Oreo Go Edition



<u>Android 9.0</u> Pie

New user interface for the quick settings menu Improved Adaptive Brightness feature A new gesture-based system interface A new "Lockdown" mode which disables biometric authentication once activated.

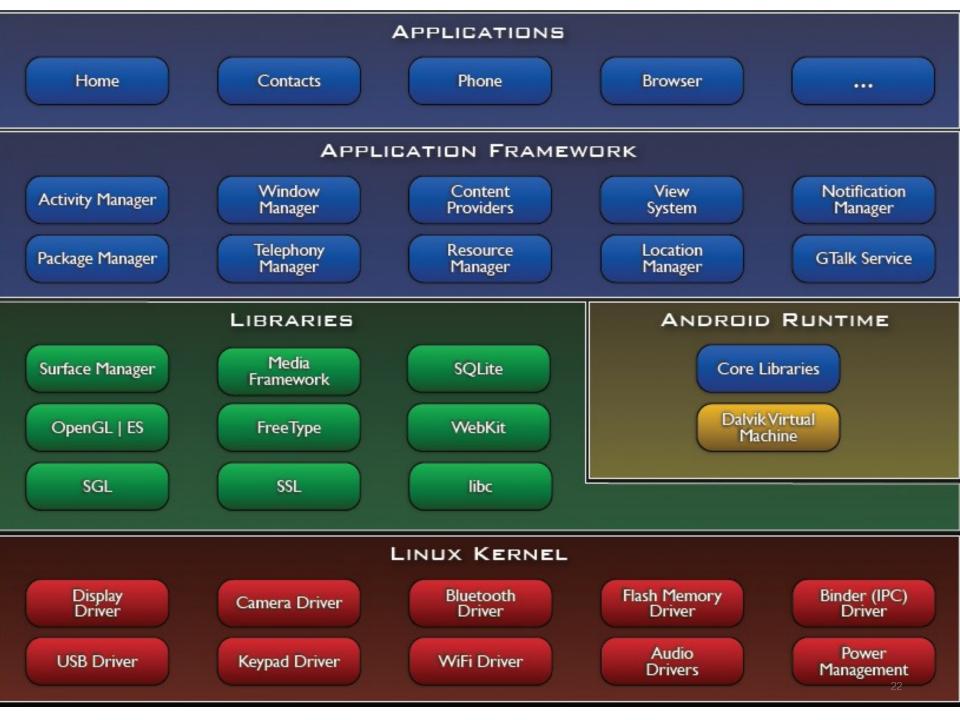
September 3, 2019

Android 10 (API 29)

- •New permissions to access location in background and to access photo, video and audio files.
- •Support for the <u>AV1</u> video codec, the <u>HDR10+</u> video format and the <u>Opus</u> audio codec.
- •A native <u>MIDI</u> API, allowing interaction with music controllers.
- •Support for the <u>WPA3</u> Wi-Fi security protocol.
- •Support for <u>foldable phones</u>.

Android Handsets





The Kernel

Why Linux kernel ?

- Great memory and process management
- Permissions-based security model
- Proven driver model
- Support for shared libraries
- It's already open source!

Standard Linux 2.6.24 Kernel

Patch of "kernel enhancements" to support Android



Android is not Linux !

- Android is built on the Linux kernel, but Android is not Linux
- No native windowing system
- No glibc support
- Does not include the full set of standard Linux utilities



Kernel Enhancements

- Alarm
- Ashmem
- Binder
- Power Management
- Low Memory Killer
- Kernel Debugger
- Logger



Power Management

- Built on top of standard Linux Power Management (PM)
- More aggressive power management policy
- Components make requests to keep the power on through "wake locks"
- Supports different types of wake locks



Libraries

Libraries

- C/C++
- Draw Pixel
- Multimedia / Codec
- Communication
- Database
- Browser
- Fonts



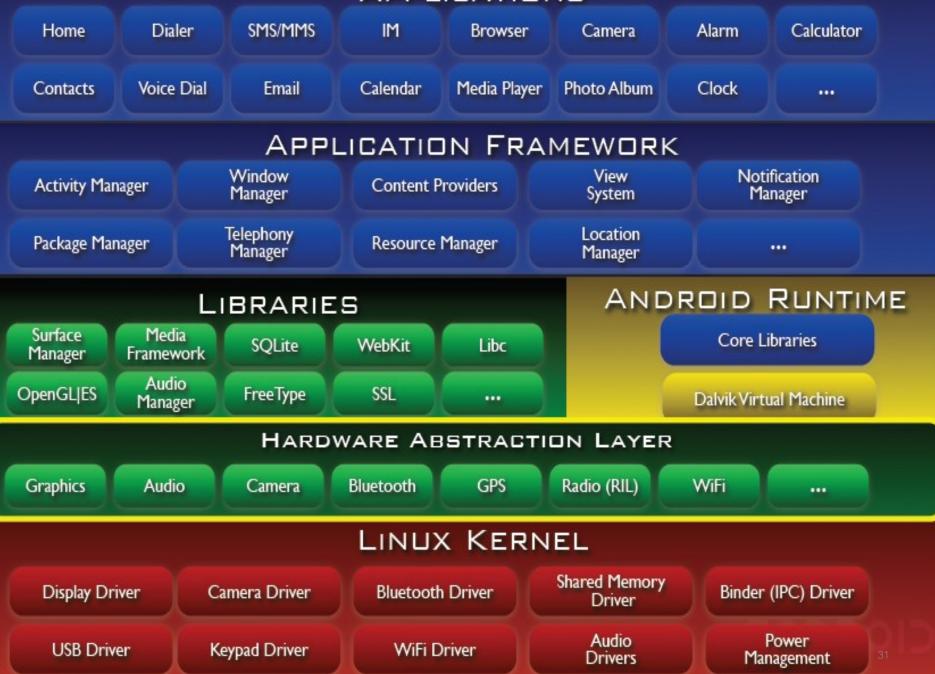
Why android using BSD Libc ?

- License: we want to keep GPL out of user-space
- Size: will load in each process, so it needs to be small
- Fast: limited CPU power means we need to be fast

X Doesn't support certain POSIX features
X Not compatible with Gnu Libc (glibc)
X All native code must be compiled against bionic



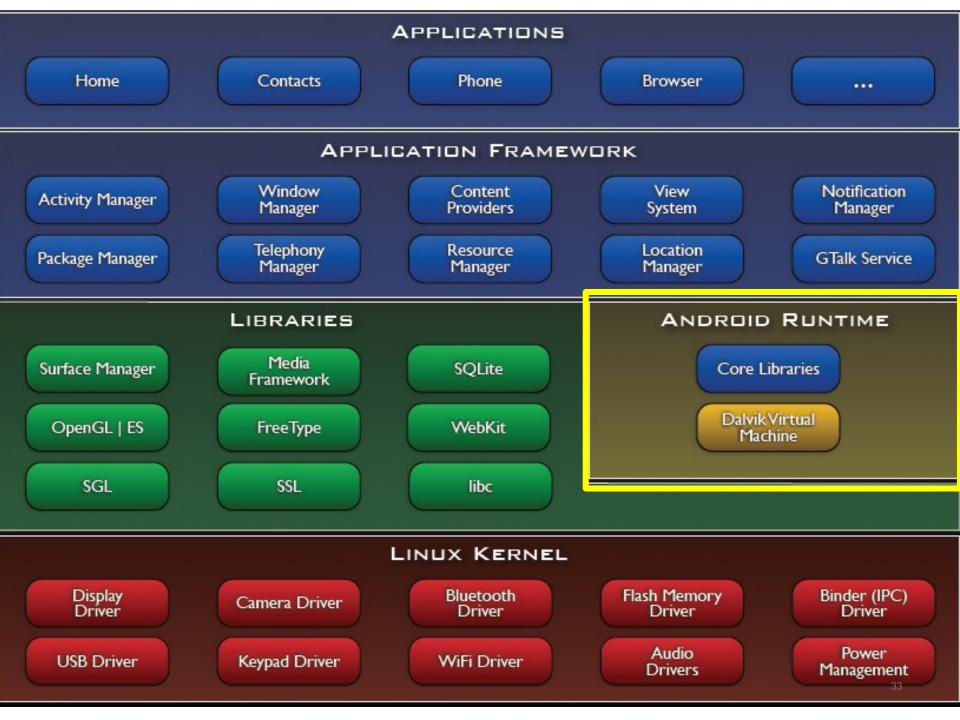
APPLICATIONS



Why do we need a user-space HAL?

- Not all components have standardized kernel driver interfaces
- Kernel drivers are GPL which exposes any proprietary IP
- Android has specific requirements for hardware drivers

HARDWARE ABSTRACTION LAYER Graphics Audio Camera Bluetooth GPS Radio (RIL) WiFi ... 32



Android Runtime Dalvik Virtual Machine



Designed for embedded environment

- Supports multiple virtual machine processes per device
- Highly CPU-optimized bytecode interpreter
- Uses runtime memory very efficiently

Application runs in sand box



```
0000: lconst 0
 0001: Istore 1
 0002: aload 0
 0003: astore 3
 0004: aload 3
 0005: arraylength
 0006: istore 04
 0008: iconst 0
 0009: istore 05
  000b: iload 05 // rl ws
  000d: iload 04 // rl ws
  000f: if_icmpge 0024 // rs rs
  0012: aload_3 // rl ws
  0013: iload 05 // rl ws
  0015: iaload // rs rs ws
  0016: istore 06 // rs wl
  0018: lload_1 // rl rl ws ws
  0019: iload 06 // rl ws
  001b: i2l // rs ws ws
  001c: ladd // rs rs rs rs ws ws
  001d: lstore_1 // rs rs wl wl
  001e: iinc 05, #+01 // rl wl
  0021: goto 000b
 0024: lload 1
 0025: Ireturn
```

```
public static long sumArray(int[] arr) {
    long sum = 0;
    for (int i : arr) {
        sum += i;
    }
    return sum;
} JavaVM
```

- 25 bytes
- 14 dispatches
- 45 reads
- 16 writes

Rs=read stack Rw=write stack Rl=read local Wl=write local

Virtual Machine Showdown: Stack Versus Registers

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ABSTRACT

Virtual machines (VMs) are commonly used to distribute programs in an architecture-neutral format, which can easily be interpreted or compiled. A long-running question in the design of VMs is whether stack architecture or register architecture can be implemented more efficiently with an interpreter. We extend existing work on comparing virtual stack and virtual register architectures in two ways. Firstly, our translation from stack to register code is much more sophisticated. The result is that we eliminate an average of more than 47% of executed VM instructions, with the register machine bytecode size only 25% larger than that of the corresponding stack bytecode. Secondly we present an implementation of a register machine in a fully standardcompliant implementation of the Java VM. We find that, on the Pentium 4, the register architecture requires an average of 32.3% less time to execute standard benchmarks if dispatch is performed using a C switch statement. Even if more efficient threaded dispatch is available (which requires labels as first class values), the reduction in running time is still approximately 26.5% for the register architecture.

Categories and Subject Descriptors

D.3 [Software]: Programming Language; D.3.4 [Programming Language]: Processor—Interpreter

General Terms

Performance, Language

Keywords

Interpreter, Virtual Machine, Register Architecture, Stack Architecture

1. MOTIVATION

Virtual machines (VMs) are commonly used to distribute programs in an architecture-neutral format, which can easily

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VEE'05, June 11-12, 2005, Chicago, Illinois, USA. Copyright 2005 ACM 1-59593-047-7/05/0006...\$5.00. M. Anton Ertl Institut für Computersprachen TU Wien Argentinierstraße 8 A-1040 Wien, Austria anton@complang.tuwien.ac.at

be interpreted or compiled. The most popular VMs, such as the Java VM, use a virtual stack architecture, rather than the register architecture that dominates in real processors.

A long-running question in the design of VMs is whether stack architecture or register architecture can be implemented more efficiently with an interpreter. On the one hand stack architectures allow smaller VM code so less code must be fetched per VM instruction executed. On the other hand, stack machines require more VM instructions for a given computation, each of which requires an expensive (usually unpredictable) indirect branch for VM instruction dispatch. Several authors have discussed the issue [12, 15, 11, 16] and presented small examples where each architecture performs better, but no general conclusions can be drawn without a larger study.

The first large-scale quantitative results on this question were presented by Davis et al. [5, 10] who translated Java VM stack code to a corresponding register machine code. A straightforward translation strategy was used, with simple compiler optimizations to eliminate instructions which become unnecessary in register format. The resulting register code required around 35% fewer executed VM instructions to perform the same computation than the stack architecture. However, the resulting register VM code was around 45% larger than the original stack code and resulted in a similar increase in bytecodes fetched. Given the high cost of unpredictable indirect branches, these results strongly suggest that register VMs can be implemented more efficiently than stack VMs using an interpreter. However, Davis et al's work did not include an implementation of the virtual register architecture, so no real running times could be presented.

This paper extends the work of Davis et al. in two respects. First, our translation from stack to register code is much more sophisticated. We use a more aggressive copy propagation approach to eliminate almost all of the stack load and store VM instructions. We also optimize constant load instructions, to eliminate common constant loads and move constant loads out of loops. The result is that an average of more than 47% of executed VM instructions are eliminated. The resulting register VM code is roughly 25% larger than the original stack code, compared with 45% for Davis et al. We find that the increased cost of fetching more VM code involves only 1.07 extra real machine loads per VM instruction eliminated. Given that VM dispatches are much more expensive than real machine loads, this indicates strongly that register VM code is likely to be much

47% of instructions were eliminate

Bytecode size only 25% increase

ACM VEE'05, June 11-12, 2005

```
0000: const-wide/16 v0, #long 0
0002: array-length v2, v8
0003: const/4 v3, #int 0
0004: move v7, v3
0005: move-wide v3, v0
0006: move v0, v7
0007: if-ge v0, v2, 0010 // r r
                                       }
0009: aget v1, v8, v0 // r r w
000b: int-to-long v5, v1 // r w w
000c: add-long/2addr v3, v5 // rrrrw w
000d: add-int/lit8 v0, v0, #int 1 // r w
000f: goto 0007
0010: return-wide v3
```

```
public static long sumArray(int[] arr) {
  long sum = 0;
  for (int i : arr) {
    sum += i;
  }
  return sum;
                          DalVikVM
                 • 18 bytes
                 • 6 dispatches
```

```
• 19 reads
```

• 6 writes

ART vs Dalvik

Desktop backup password Desktop full backups aren't currently protected Select runtime Use Dalvik

Use ART

Cancel
Process Stats
Geeky stats about running processes Dalvik is based on **JIT (just in time) compilation**.

ART, on the other hand, compiles the intermediate language, Dalvik bytecode, into a **system-dependent binary**.

Application Framework & Toolkit

Toolkit for android application

- Activity Manager
- Resource Manager
- Content Provider
- Notification Manager



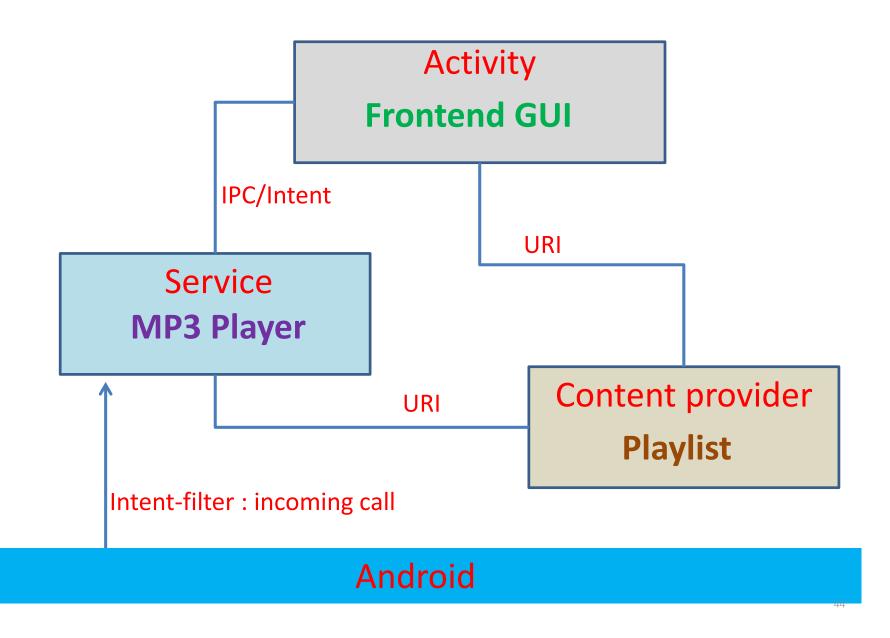
Android Application

Application Building Blocks

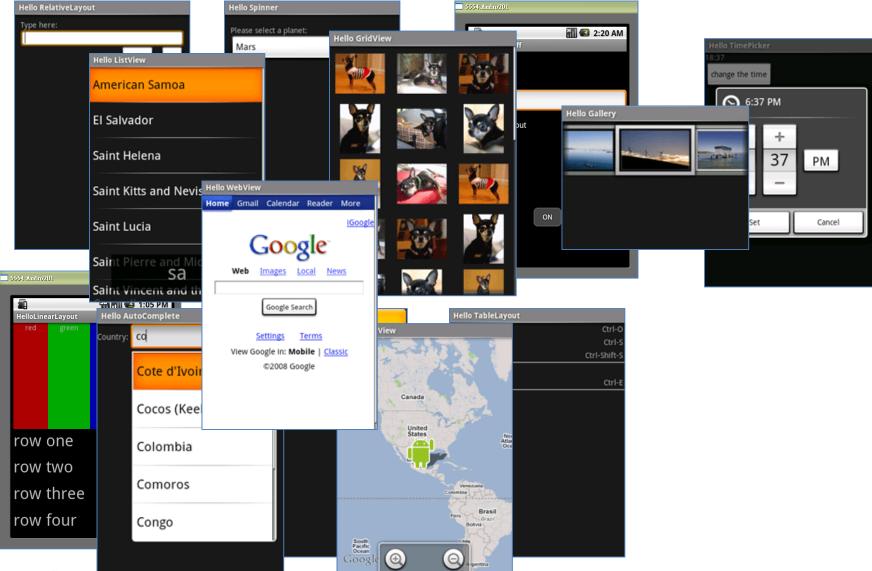
Activity :UI component typically corresponding to one screen. IntentReceiver: Set and respond to notifications or status changes. Can wake up your app. Service: Faceless task that runs in the background. ContentProvider: Enable applications to share data.



Android application : Music Player



User Interface



Android Apps



Sleep as android



Run keeper

Android Apps

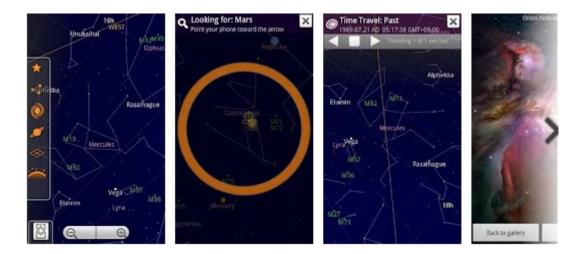


Shazam



gStrings

Android Apps

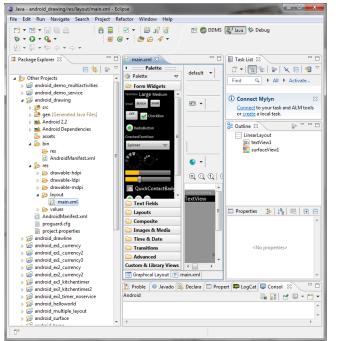


Google skymap



Games

Android software development





-J2se JDK -Android SDK -SDK Platform -SDK Platform tools -SDK Tools -Emulator & Images -Example -Eclipse IDE

-ADT Plug-in

-Android Studio

-ADB USB Driver

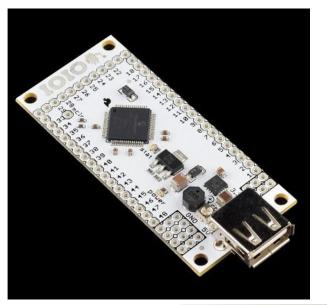
-Internet connection for online install

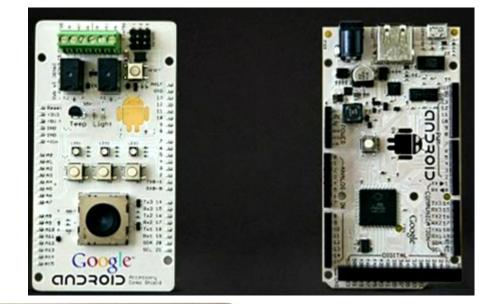
Emulator limitation

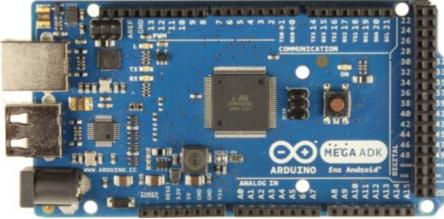


No support for placing or receiving actual phone calls. No support for USB connections No support for camera/video capture (input). No support for device-attached headphones No support for determining connected state No support for determining battery charge level No support for determining SD card insertion/removal No support for Bluetooth No support for Multitouch

Android Open Accessory Development Kit (ADK)







IOIO for Android



QA